Claims

We claim:

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5UB	A method for maintaining a system for database management, the method comprising
2	during splitting of a leaf block of a database index recording an address of a newly
3	created leaf block; and
4	maintaining the new address in a list as part of metadata of a primary B+tree.
1	2. The method according to claim 1, further comprising:
	maintaining a measure of invalid guess-database block addresses by calculating a ratio of
	a count of database block addresses in the list of new addresses to a total number of leaf blocks
4	of the primary B+tree.
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W W	3. The method according to claim 2, wherein the measure of invalid guess-database
	block addresses applies to mapping tables and secondary indexes on the primary B+tree.

- block addresses applies to mapping tables and secondary indexes on the primary B+tree.
- 4. The method according to claim 2, wherein the list of database block addresses and the ratio are maintained only when the ratio is less than a threshold value.
- 5. The method according to claim 4, wherein the threshold value for the ratio is about 1 2 10%.

1	6. The method according to claim 3, further comprising:
2	adjusting a guess-DBA quality of at least one of the mapping table and the secondary
3	index utilizing the ratio.
1	7. The method according to claim 4, wherein if the ratio is below the threshold value the
2	method further comprises:
3	selectively correcting entries in the mapping table and/or secondary index.
1	8. The method according to claim 7, wherein correcting entries in the mapping table
(<u>1</u> 2	comprises for all rows in a list of blocks in the primary B+tree:
A-SECTO	obtaining corresponding mapping table row identifiers and database block addresses of a
14	current block in the list;
	sorting the corresponding mapping table row identifiers;
<u>.</u>	obtaining mapping table rows corresponding to the mapping table row identifiers; and
	updating a guess-DBA component if it has changed.
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1	9. The method according to claim 8, wherein the correcting is carried out on-line in a
2	piece-wise manner.
1	10. The method according to claim 4, wherein correcting entries in the secondary index
2	comprises for all rows in a list of blocks in the primary B+tree:
3	obtaining a secondary index key, a primary key and a database block address of a current
4	block in the list of blocks;

,	sorting the secondary fildex keys, primary keys and database addresses in order of
6	(secondary index key, primary key) pairs;
7	obtaining an index row corresponding to the (secondary index key, primary key) pair; and
8	updating a guess-DBA component of the index row if the guess-DBA has changed.
1	11. The method according to claim 11, wherein the correcting is carried out on-line in a
2	piece-wise manner.
1	12. The method according to claim 4, wherein if the ratio is above the threshold value
2	the method further comprises:
	correcting guess-database addresses on a per object basis.
j	13. The method according to claim 12, wherein correcting guess-database block
2	addresses on the mapping table comprises:
	performing a full scan of the mapping table;
4	determining for each row of the mapping table a correct guess-database block address by
5	traversing the primary B+tree up to a penultimate level;
6	updating each row of the mapping table with the correct guess-database block address;
7	and
8	committing the correct guess-database address to the mapping table in batches.
1	14. The method according to claim 12, wherein correcting guess-database block
2.	addresses on a per object basis comprises for each secondary index object:

3	performing a full scan of the secondary index object;
4	determining for each row of the secondary index a correct guess-database block address
5	by traversing the primary B+tree up to a penultimate level;
6	updating each row of the secondary index with the correct guess-database block address;
7	and
8	committing the correct guess-database block address to the secondary index in batches.
1	15. The method according to claim 1, further comprising:
2	maintaining a list of database block addresses in the list.
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13	16. A system for organizing a database index, the system comprising:
	a list of addresses of blocks newly created during splitting of a primary B+tree.
() ()	17. The system according to claim 16, further comprising:
	a count of database block addresses in the list.
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1	18. The system according to claim 16, further comprising.
2	a ratio of count of database block addresses to total number of leaf blocks as a measure of
3	invalid guess-database block addresses.
1	19. The system according to claim 16, wherein the database index is a primary B+tree
2	structure, wherein the system further comprises:
3	a mapping table used to support bitmap indexes.

1	\20. The system according to claim 19, further comprising:
2	a bitmap index supported by the mapping table.
1	21. The system according to claim 16, wherein the database index is a primary B+tree
2	structure, wherein the system further comprises:
3	a secondary index structure comprising hybrid row identifiers.
1	22. A computer program product for performing a process for maintaining a database
	management system, comprising:
3	a computer readable medium; and
4	computer program instructions, recorded on the computer readable medium, executable
5	by a processor, for performing the steps of:
б	during splitting of a leaf block of a primary B+tree recording an address of a newly
6 7 7 8	created leaf block; and
: 6	maintaining the new address in a list as part of primary B+tree metadata.
1	23. A system for performing a process for maintaining a database management system,
2	comprising:
3	a processor operable to execute computer program instructions; and
4	a memory operable to store computer program instructions executable by the processor,
5	for performing the steps of:
6	during splitting of a leaf block of a primary B+tree recording an address of a newly



- 7 created leaf block; and
- 8 maintaining the new address in a list as part of primary B+tree metadata.